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(71) Applicant: XEROX CORPORATION  
Rochester, New York 14644 (US)

(72) Inventors:

- Nguyen, Uoc H  
Long Beach, California 90807 (US)
- Claproth, Abraham E.  
Culver City, California 90230 (US)
- Nguyen, Kien T.  
Torrance, California 90501 (US)

- Lee, Chio-hao  
Hacienda Heights, California 92745 (US)
- De Queiroz, Ricardo L.  
Pittsford, NY 14534 (US)
- Lao, Guillermo  
Torrance, California 90503 (US)
- Romano, Juan A.  
North Hills, California 91343 (US)
- Tsai, Tim S.  
Irvine, California 92620 (US)

(74) Representative:

Skone James, Robert Edmund  
GILL JENNINGS & EVERY  
Broadgate House  
7 Eldon Street  
London EC2M 7LH (GB)

### (54) HVQ compression with an error term

(57) Improving the accuracy of a lossy compressor (10) by decompressing the compressed input word and comparing the result to the original word. An error term is produced, and the total compressor output becomes the output of the lossy compressor and the error term.

The compression can be improved by limiting the error term to a number of most significant bits, and by losslessly compressing both terms. An HVQ compressor can be used as the lossy compressor.

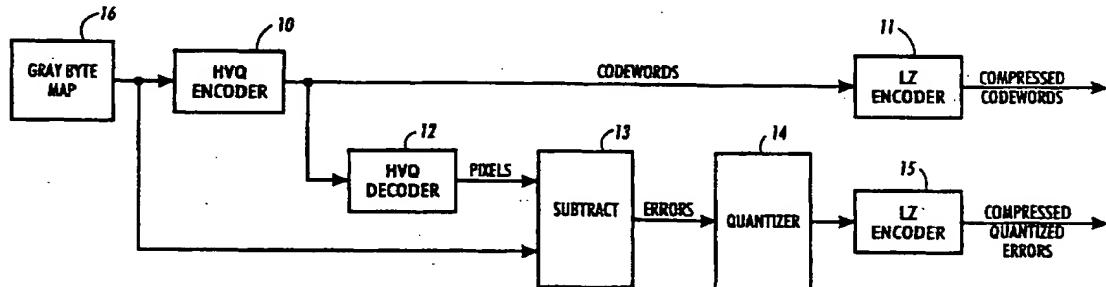


FIG. 1

## Description

[0001] When data must be transmitted or stored it is usually compressed first, to reduce the transmission time or storage requirement. This is especially true of image data which can consist of four color separations of 8-bit pixels.

[0002] One method of compression is HVQ, where a block of data can be reduced to a single codeword in a number of steps. This method is described in U.S. Patent Number 5,602,589, and is incorporated by reference herein. This procedure can be explained most clearly by the use of a numerical example in connection with Figure 1a of that patent. Assume that a 2 by 4 pixel block, 8 bits per pixel, is to be compressed to one 9-bit number. First, each pair of pixels is applied to a stage 1 look up table (LUT) containing 9-bit code words, each code word being associated with a two-pixel bit pattern. The table is set up so that if an exact match of the input pixel bits is not available, then a codeword associated with the closest match is output instead. The ultimate output of this stage 1 is four codewords, each describing the closest match to the bit pattern of the two input pixels. Since the exact match frequently is not possible, this compression is inherently lossy to some extent. The output of the first stage is four 9-bit codewords.

[0003] Assuming this amount of compression is not sufficient, these four codewords, each representing the bit pattern of a 1 by 2 pixel block are next applied to two stage 2 tables, resulting in a total of two 9-bit output codewords, each representing the bit pattern of a 2 by 2 pixel block. If the compression is still insufficient, these two are applied to the stage 3 table to yield the final output word, which is associated with a 2 by 4 pixel block. It can now be seen that this method is hierarchical in that a plurality of stages are used, it involves vectors since the input pixels have a directional relationship with each other, and is quantized in that a number of levels can be implemented to achieve any degree of compression.

[0004] For decompression, each codeword is simply applied to a 256 K by 64 bit LUT which outputs the bit pattern of the entire 8-pixel block.

[0005] In accordance with one aspect of the present invention, a process of compressing an original multi-bit word comprises the steps of compressing the original word using a lossy compressor to create a compressed word which comprises one part of the output,

decompressing the compressed word to form a decompressed word,

comparing the decompressed word and the original word to form an error word which forms the other part of the output.

[0006] In accordance with a second aspect of the present invention, a data compression system comprises a lossy compressor for compressing an original

multi-bit word to create a compressed word which comprises one part of the output, a decompressor for decompressing the compressed word to form a decompressed word; and a comparator for comparing the decompressed word and the original word to form an error word which forms the other part of the output.

[0007] An additional element that can be added to the general HVQ system to improve its performance is an error channel. After compression, the compressor decompresses the final output codeword to form the associated bit pattern, and then compares that to the original input pixel block bit pattern. If there was any loss in the compression process, then an error can be calculated. This error can then be sent along with the compressed output to the decompressor. After the decompressor has produced the final bit pattern, the error term can be added to it to produce an errorless final product. Of course, if the entire error term is sent along with the compressed output, the compression ratio will suffer. The best tradeoff is to send only a few of the most significant bits of the error term, perhaps only two or three bits. In this way, the final result is improved accuracy and/or compression.

[0008] The performance of a "hierarchical vector quantization" (HVQ) circuit can thus be improved by using the compressor to decompress the data, compare the result to the original data to produce an error term, and send the error term to the decompressor to correct errors in the final decompressed data.

[0009] Some examples of methods and system according to the present invention will now be described with reference to the accompanying drawings, in which:-

Figure 1 is an HVQ encoder which has a parallel error channel;  
 Figure 2 is the decoder for the arrangement of Figure 1;  
 Figure 3 is a one-channel arrangement for coding a pixel split into its most and least significant bits;  
 Figure 4 is a one-channel arrangement for decoding a pixel split into its most and least significant bits;  
 Figure 5 is a two-channel arrangement for coding a pixel split into its most and least significant bits;  
 Figure 6 is a decoder for the coder of Figure 5;  
 Figure 7 shows three types of rotating or mirror imaging an image using HVQ compression;  
 Figure 8 shows a boundary between blocks of different sizes; and,  
 Figure 9 shows a boundary between blocks of the same size.

[0010] The basic HVQ system can be improved by adding an error channel as shown in Figure 1. In the upper channel the grayscale byte map 16 is applied in the usual way to an HVQ decoder 10, the output is losslessly compressed in an LZ encoder 11, and the result is sent to the decoder, usually in the form of 8 to 10-bit

words.

[0011] In addition, the output of the HVQ encoder is sent to a decoder 12 in scanline format which produces a version of the original byte map which may be different from the original because of errors possibly introduced by the lossy encoder. The two byte maps are then subtracted 13, pixel by pixel, to produce error terms which, if added to the output codeword, will produce the original byte map. This subtraction can also be done by using an exclusive OR, which is simpler and faster and does not require a sign bit. These error terms, each a signed quantity 8-bits wide or less, can then be compressed in an LZ encoder 15 and sent to the decoder in parallel with the original output. The larger the error term, the less will be the compression ratio. In practice, small errors are not visually detectable. To limit the degradation of the compression ratio, the error term can be limited to a few most significant bits, three for example, in quantizer 14. Normally, the amount of error for a pixel will not be large enough to show up in the few MSB's, in which case there will be no error term at all.

[0012] The quantized error decoder is shown in Figure 2. The compressed code words are LZ decoded in decoder 21 and HVQ decoded in decoder 22 to produce one term for the adder 23. The compressed quantized errors are LZ decoded in decoder 24 and applied as the other term to the adder 23, the output of which is output video. The adder 23 can either be an adder adding a sign bit and seven bits, or an exclusive OR if one was used to generate the error term in the encoder.

[0013] Figure 3 shows the arrangement when a single code word 31 is split into most and least significant parts, 32, 33, and where only the least significant bits are compressed. In this case bits 0 through 4 are sent through lossy compressor 35 while bits 5 through 7 are not. Both are then compressed using lossless LZ compression 34,36 and output to the decoder shown in Figure 4. Here again, both channels are LZ decompressed 41,42 while only the LSB's are HVQ decoded 43. The two resultant parts are then exclusive ORed 44 together to be applied to the decoding look up table.

[0014] Figure 5 is a system where a single pixel is separated into a least significant segment and a most significant segment, and where a separate and different compression process 54,56 is used for each segment, the least significant bits being more compressed. The original pixel is separated into its most significant bits 52 and least significant bits 53. The result is that the most significant bits, being the most important, are less compressed while compression for the least significant bits has a better compression ratio. A programmable look-up table could be used to split the input pixel into any two segments other than the 3 - 5 split shown.

[0015] Figure 6 is the decoder for the encoder of Figure 5. The two compressed outputs of Figure 5 are applied to LZ decoders 61, 62 and HVQ 63, 64 decoded. Then both are applied to an exclusive OR gate 65 to assemble the entire pixel. Of course, if the

encoder used some other combination of encoders, the decoder would use the same form of decoding. That is, more generally, the data words in a string can be divided into more and less significant bits to create two parallel strings, and then compressed using any two methods of compression where the greater compression is applied to less significant bits.

[0016] HVQ compression is easily adapted to image rotation and mirror imaging as shown in Figure 7. The process is shown here using the example of an original image that is four pixels high and sixteen pixels wide, and must be rotated ninety degrees clockwise and mirror imaged.

[0017] Step 1 is the usual compression process of reducing the eight 8-pixel segments to eight codewords Cw1 to Cw8. Step 2 is the step of rearranging the codewords into the rotated and mirror imaged order. This hardware can be in the form of wiring where the second word in, for example Cw2, is connected to the third word out, as shown. Step 2 can have several sets of wiring, each set delivering a different rotation. Finally, decoding step 3 uses a look-up table to produce a pixel pattern for each segment that is properly oriented. Here again, several tables can be used to produce various orientations.

[0018] Printing hints may be incorporated into the original data supplied by the user in the original page description language to indicate to the printer how the data may best be printed. For example, a hint word may be two bits in length, and indicate one of four possibilities, that the following data is text, contone, graphics, etc. For example, if the printer is receiving data that originated as a computer generated graphic it may use a different halftone screen than it would if the original data was screened in from a photograph.

[0019] Printing hints may be added to any HVQ channel as shown in Figure 5. Assume that for each 4 pixel block entering into the HVQ encoder 54 there is produced one codeword 9 bits in length, contained in two 8-bit bytes, so that the first 8 bits are contained in the first byte and the last bit is contained in the second byte. Then the 2-bit hint is added. Now, each codeword plus hint is 11 bits, still contained in two bytes. The LZ encoder 55 looks at a string of bytes, perhaps 256 bytes in length, and determines the location and size of the most recent identical pattern match. To the extent that the hint changes once or twice within that string, there will be a slightly decreased amount of compression in comparison to the amount that would have resulted with no hints. However, to the extent that the hint does not change during that time, there is no decrease in compression at all. The result is that printing hints can be supplied after an HVQ compressor but before the lossless compressor with very little effect on the compression ratio. At the output side, after the code words are LZ decompressed, but before being decoded, the hints can be extracted for later use.

[0020] The losses of an HVQ compressor can be fur-

ther minimized by choosing codewords and output data patterns that have the best chance of matching the actual input data patterns. For example, first consider text. If text pixels are being encoded in 4 by 2 pixel groups, and the four input pixels in one line are black, dark gray, light gray and white, and the input video was scanned-in text, the data most likely originally was a boundary between a black letter and a white space, so the output pixel pattern could be black, black, white, white. On the other hand, if the original input data was a scanned-in computer generated graphic, the four pixels are more likely to be a smooth decrease in density from black to white. The actual determination of the encoder codewords and patterns in the decoder look-up table are determined by statistical analysis. A representative group of text and graphic documents are passed through a test program and the best values are generated for each type.

[0021] A complication arises when a boundary passes through an input block of pixels, in which case neither text nor graphic values can be used for the entire block. The solution is to supply a third set of codewords and patterns which are generated specifically for this boundary condition. In this case a set of documents containing both text and graphics would be analyzed to produce one set of patterns, which would be used when a boundary, mixed, condition is determined to be within the block.

[0022] The boundary condition is sensed by observing the printing hints. For example, a rectangular scanned-in picture is typically located on a page of text by its x,y coordinates. When the raster output scanner beam is within the coordinates, the printing hints will indicate to the printer which codewords, look-up table entries and halftone screens to use. If the hint changes from picture to text within the block, for example, then the encoder knows that a boundary exists within the block.

[0023] In all cases the block size must be maintained. Figure 8 is an example of a transition between text and contone. Since different block sizes can be used in HVQ encoders for different kinds of data, the block size for the text is shown here as 2 by 2 pixels to allow greater edge detail while the block size for contone is shown as being 4 by 2 to allow for greater compression.

[0024] If a boundary is within a 2 by 2 pixel block, that block is encoded and decoded using boundary values. In addition, any contone pixels to the right of the boundary such as pixel 81 are also treated as a boundary pixels if necessary so that all remaining pixels to the right of the boundary line will be within 4 by 2 pixel blocks. Similarly, in Figure 9, showing a boundary between 4 by 2 pixel blocks of text and contone, 2 by 2 pixel boundary blocks are used in pairs so that all remaining blocks will be 4 by 2.

## Claims

## 1. The process of compressing an original multi-bit

word comprising the steps of compressing the original word using a lossy compressor (10) to create a compressed word which comprises one part of the output,

decompressing the compressed word to form a decompressed word.

comparing the decompressed word and the original word to form an error word which forms the other part of the output.

2. The process of claim 1, further comprising the step of losslessly compressing each part of the output.

15. 3. The process of claim 1, wherein only the most significant bits of the error word are output.

4. The process of any of the preceding claims, wherein the lossy compressor (10) is an HVQ compressor.

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5. A data compressing system comprising a lossy compressor (10) for compressing an original multi-bit word to create a compressed word which comprises one part of the output; a decompressor (12) for decompressing the compressed word to form a decompressed word; and a comparator (13) for comparing the decompressed word and the original word to form an error word which forms the other part of the output.

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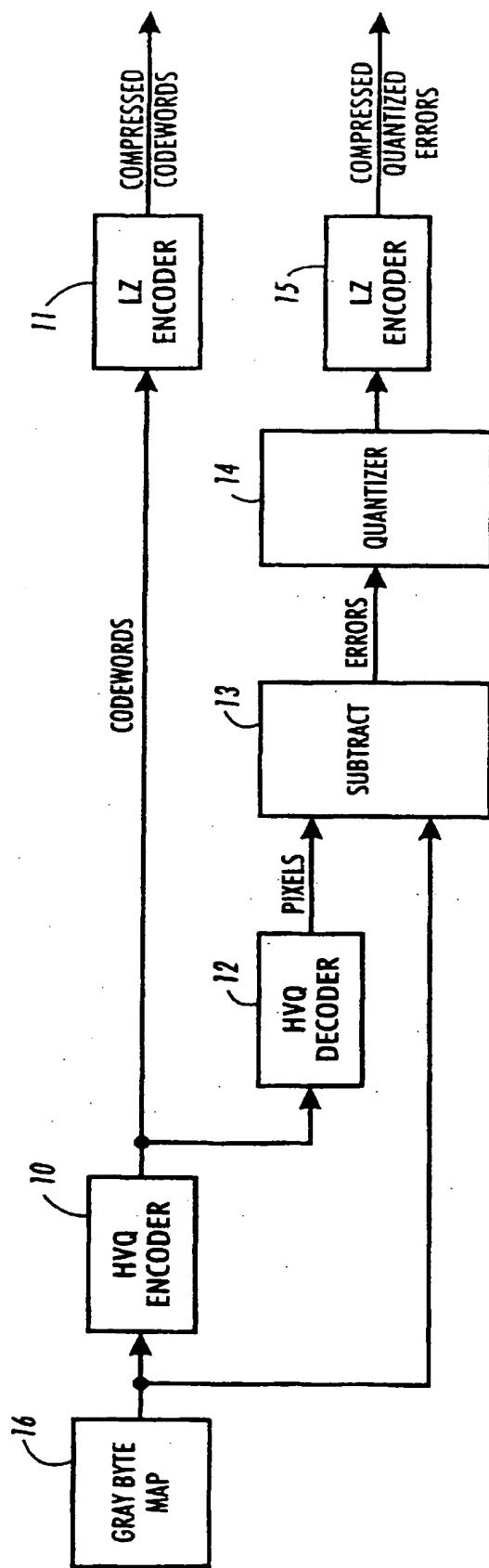


FIG. 1

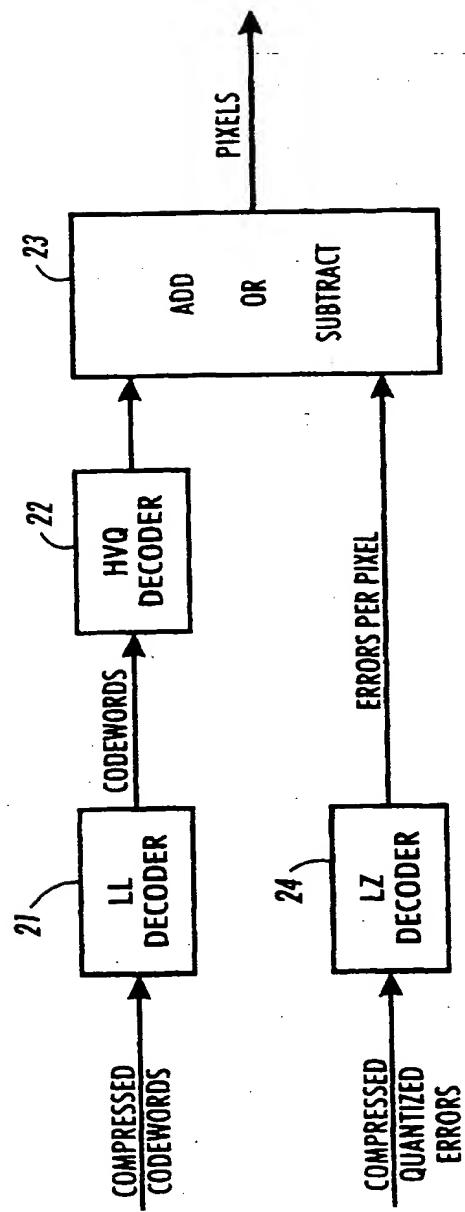


FIG. 2

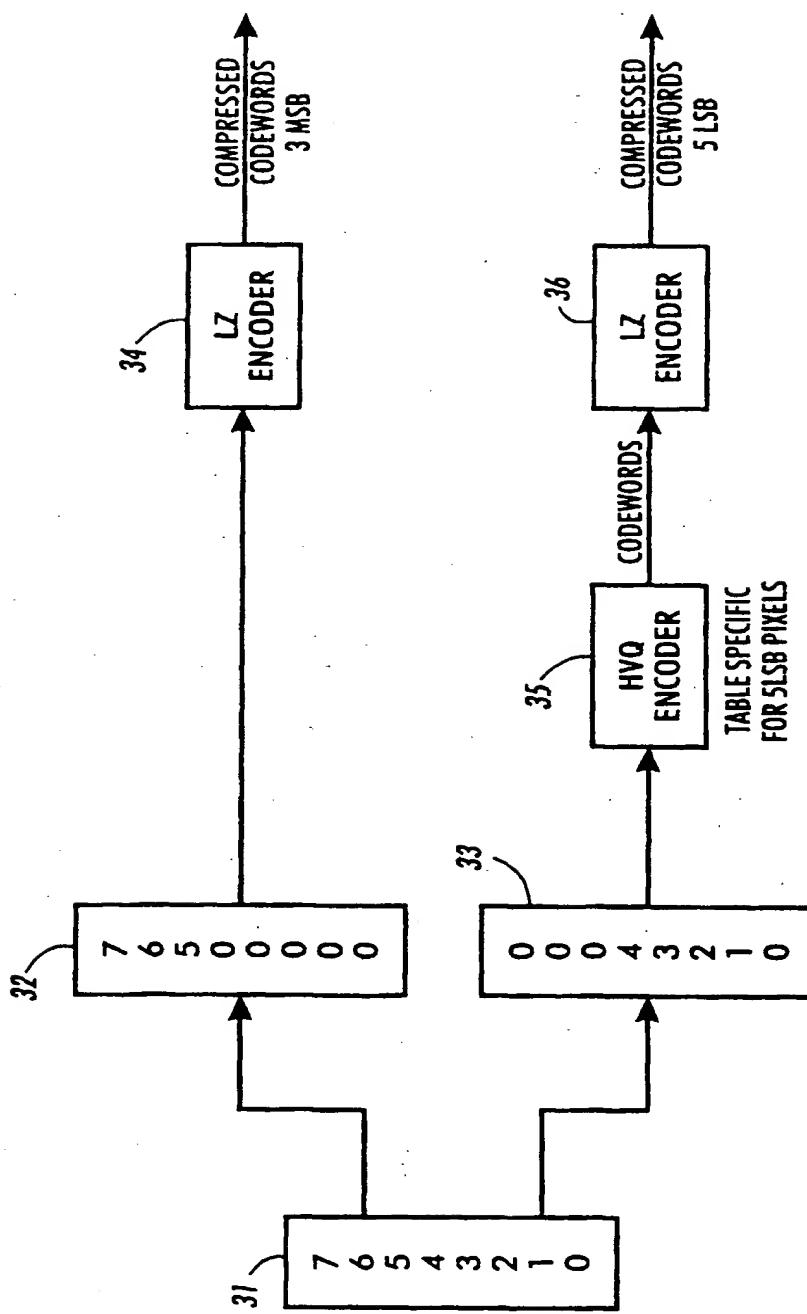


FIG. 3

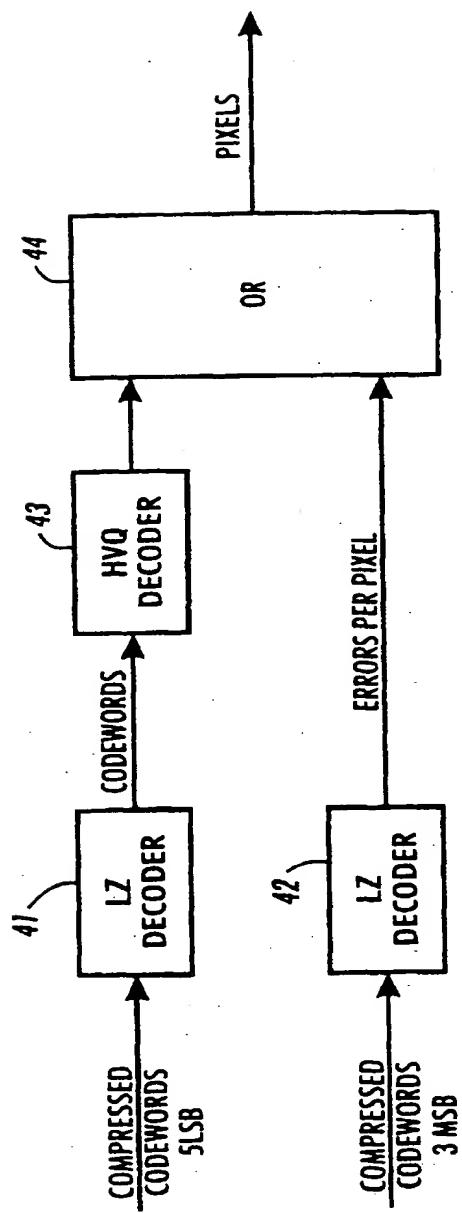


FIG. 4

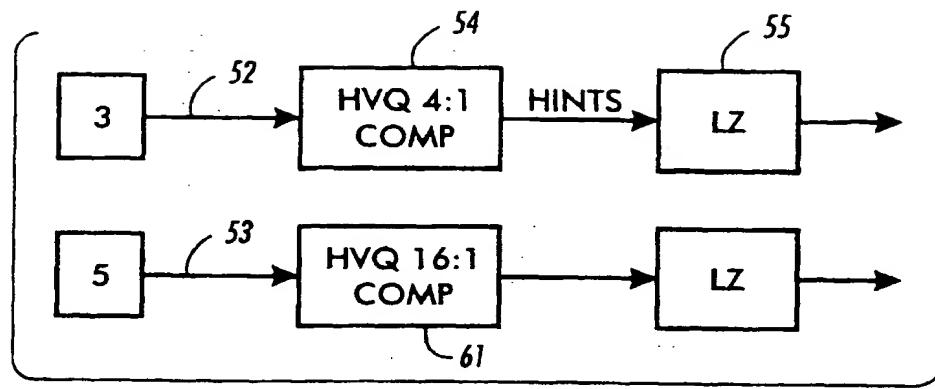


FIG. 5

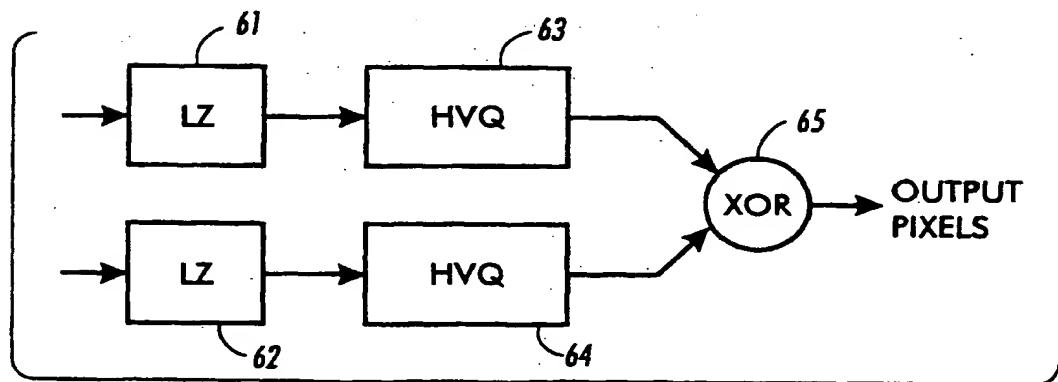
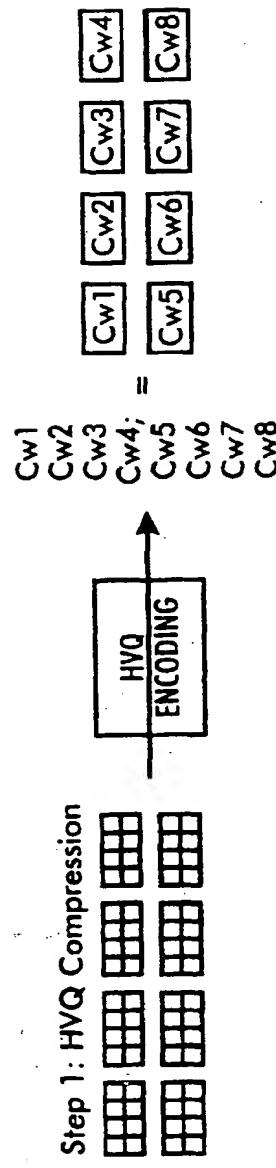
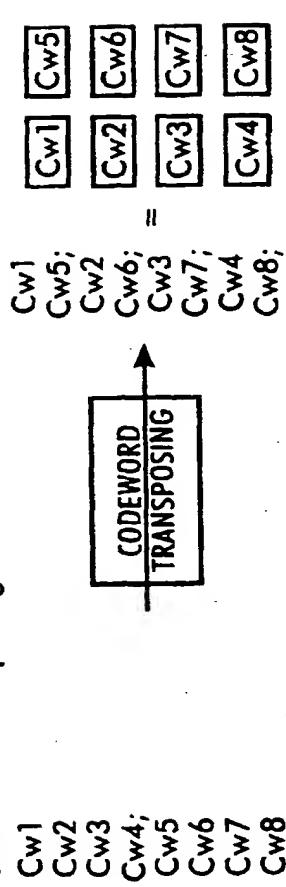


FIG. 6



## Step 2: Codeword Transposing



## Step 3: HVQ Decoding with pre-rotated tables

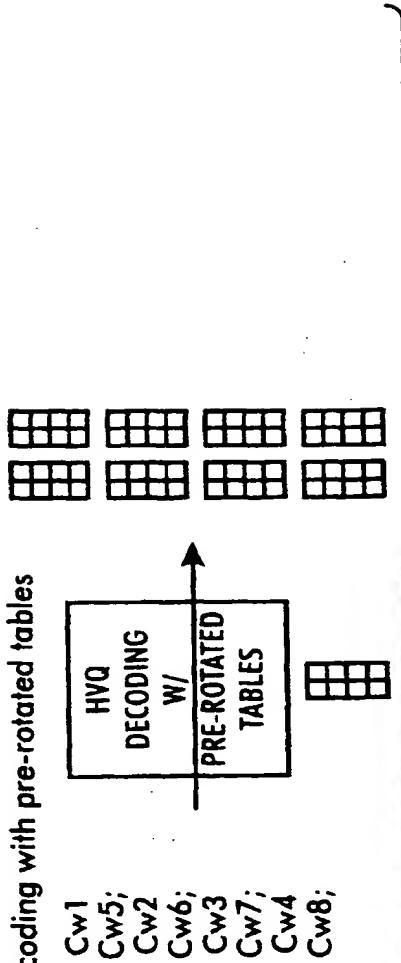
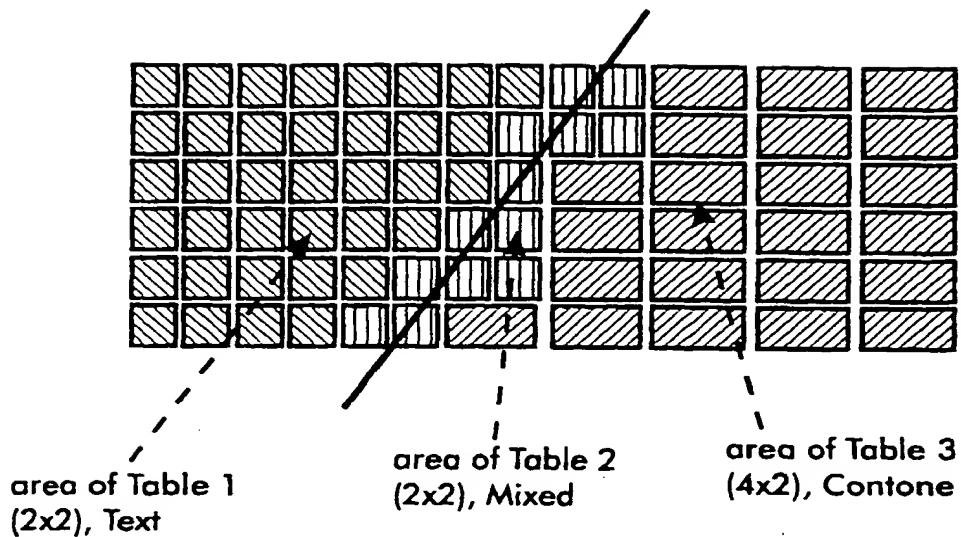
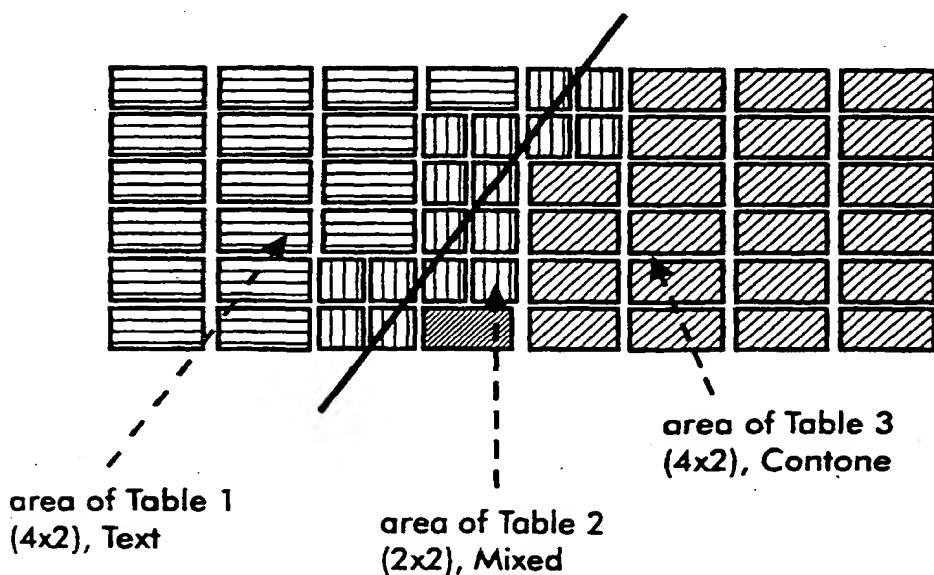


FIG. 7



**FIG. 8**



**FIG. 9**

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(71) Applicant: XEROX CORPORATION  
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(72) Inventors:

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Long Beach, California 90807 (US)
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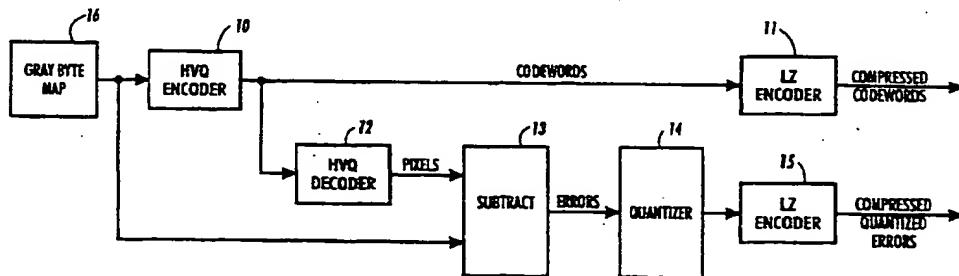


FIG. 1



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<p>The present search report has been drawn up for all claims</p> <table border="1" style="width: 100%; border-collapse: collapse;"> <tr> <td style="width: 33%;">Place of search</td> <td style="width: 33%;">Date of completion of the search</td> <td style="width: 33%;">Examiner</td> </tr> <tr> <td>THE HAGUE</td> <td>6 February 2001</td> <td>Georgiou, G</td> </tr> </table>				Place of search	Date of completion of the search	Examiner	THE HAGUE	6 February 2001	Georgiou, G
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